



#### nD-PointCloud, what else



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Congres AHN & Beeldmateriaal, 16 October 2025, Amersfoort

#### Motivation nD-PointCloud

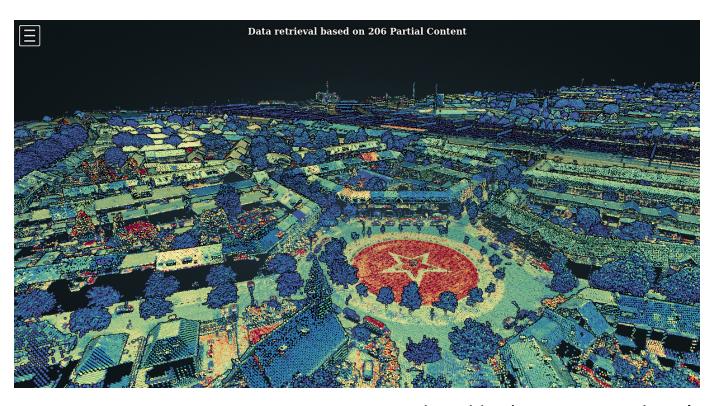
- point cloud data sets are often used for monitoring
  - → dynamic point clouds
  - time added as additional organizing dimension

- organizing point cloud data in levels of Importance (LoI) is an approach to manage large data sets
  - → LoI: discrete (multi-scale/dLoI) or continuous (vario-scale/cLoI)
  - → scale treated as additional organizing dimension

• how to manage higher dimensional point clouds (4D, 5D, ...)?

## Agenda

- nD-PointCloud
  - foundations
  - PostgreSQL implementation
  - nD Convex polytope query
  - Apache Parquet
- practical results
  - Potree conversions AHN 1 to 5
  - COPC/VPC conversions
  - fast direct point cloud-based change detection
- future work
  - main publications



AHN3 colored by 'intensity gradient'

#### nD-PC foundations

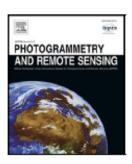
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Contents lists available at ScienceDirect

#### ISPRS Journal of Photogrammetry and Remote Sensing







#### Organizing and visualizing point clouds with continuous levels of detail

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#### ARTICLE INFO

Keywords: nD point clouds Continuous level of detail (cLoD) Space Filling Curve (SFC) Perspective view selection

#### ABSTRACT

Point clouds contain high detail and high accuracy geometry representation of the scanned Earth surface parts. To manage the huge amount of data, the point clouds are traditionally organized on location and map-scale; e.g. in an octree structure, where top-levels of the tree contain few points suitable for small scale overviews and lower levels of the tree contain more points suitable for large scale detailed views. The drawback of this solution is that it is based on discrete levels, causing visual artifacts in the form of data density shocks when creating

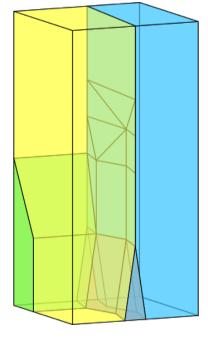
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c Alibaba Group, China

### Vario-scale for point cloud data

 lesson from vario-scale research: add one continuous dimension to the vector geometry to represent scale -> continuous Lol (2D data vario-scale represented by 3D geometry)



- apply this to point cloud data
  - 1. compute the cLol value
  - 2. add this as organizing dimension, either x,y,imp (z and others attributes) or x,y,z,imp (and others as attributes) or ...
  - 3. Cluster/index the 3D, 4D, .. nD points
  - 4. Define perspective view selections, view frustum with one more dimension: the further, the higher cLol's

## CLoI computation, getting rid of discrete levels → real continuous levels, nD case

• for ideal continuous distribution function over levels (nD):

$$f(l,n) = \frac{2^{(n-1)l}(n-1)\ln 2}{2^{(n-1)(L+1)}-1} \qquad \text{for } l \text{ between } \theta \text{ and } L+1 \\ \text{and } n \text{ number of dimensions}$$

• this function has Cumulative Distribution Function (CDF):

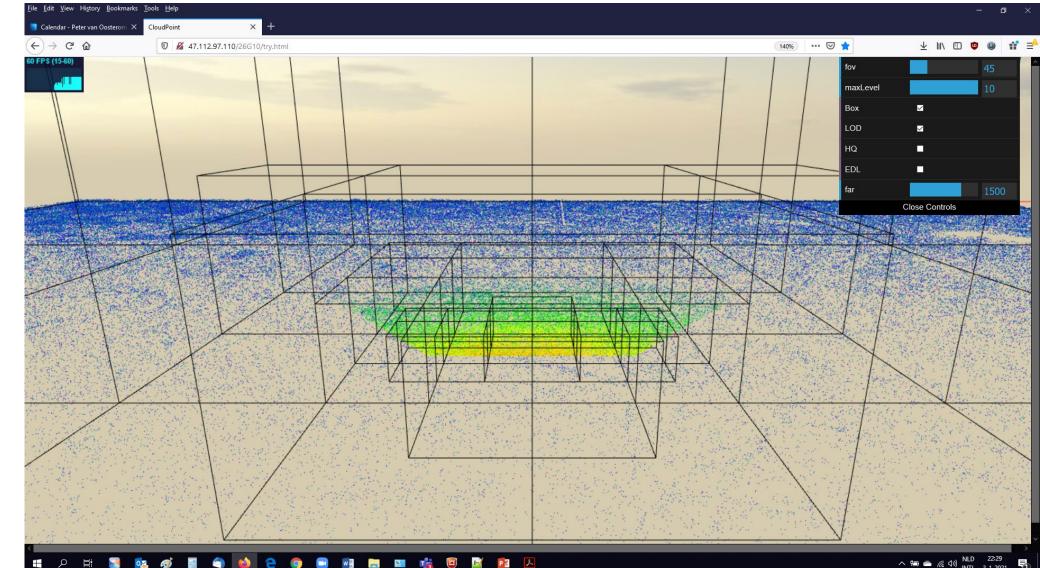
$$F(l,n) = \frac{2^{(n-1)l}-1}{2^{(n-1)(L+1)}-1}$$
 for  $l$  between  $0$  and  $L+1$  and  $n$  number of dimensions

• using random generator U (uniform between 0 and 1) to generate level l (cLoI) (between  $\theta$  and L+1) for next point in nD space:

$$l = \frac{\ln((2^{(n-1)(L+1)} - 1)U + 1)}{(n-1)\ln 2}$$

#### cLol in 3D web application

developed by Xuefeng Guan (Wuhan Univ) after 1 year visit TUD

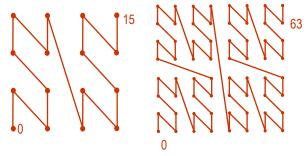


note: colour is density (3<sup>rd</sup> person view)

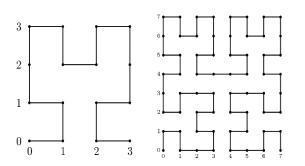
## Cluster/index by Space Filling Curves (SFCs)

- from nD to 1D (computer memory addresses) and back...
- apply linear ordering to a multidimensional domain (spatial clustering)
- organize a flat table efficiently
- full resolution keys: avoid storing x,y[,z] + t/cLoI
   → recovered from SFC key
- use Index Organized Table (cluster data stored in the B-Tree index)
- queries need to be re-written to SFC-ranges, benefit from spatial clustering → efficient
- SFCs based on hyper-cubes
  - Morton/Hilbert both nD and quadrant recursive
  - Consider relative scaling of dimensions
  - Space reserved on the hypercube for future data

Morton (Peano)



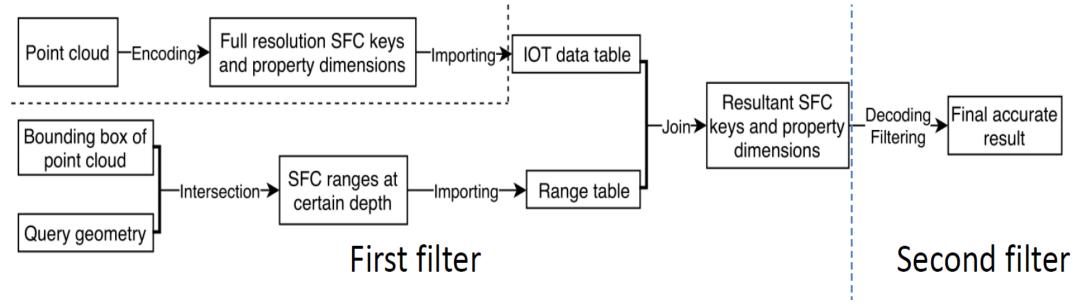
Hilbert



#### Overview: load & query point clouds

#### 1. loading (upper right of figure)

- convert nD points to SFC keys
- store in IOT table (organized on SFC key)



#### 2. querying (remainder)

- convert query geometry to SFC ranges,
- first filter: join SFC range table with IOT table
- second filter (optional): test points inside query geometry

# PostgreSQL Implementation

ISPRS Annals of the Photogrammetry, Remote Sensing and Spatial Information Sciences, Volume X-4/W2-2022 17th 3D Geolnfo Conference, 19-21 October 2022, Sydney, Australia

## PCSERVE - ND-POINTCLOUDS RETRIEVAL OVER THE WEB

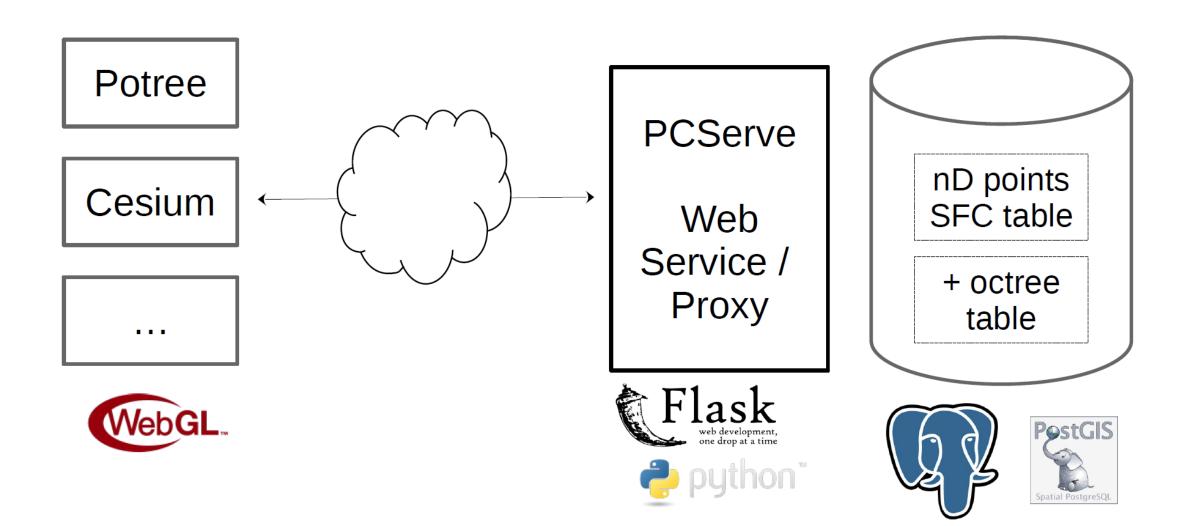
Delft University of Technology, Faculty of Architecture and the Built Environment, GIS-technology – b.m.meijers@tudelft.nl

#### Commission IV, WG IV/9

KEY WORDS: Point Clouds, Space Filling Curve, Web Based Visualisation, Web Service, Database Management System

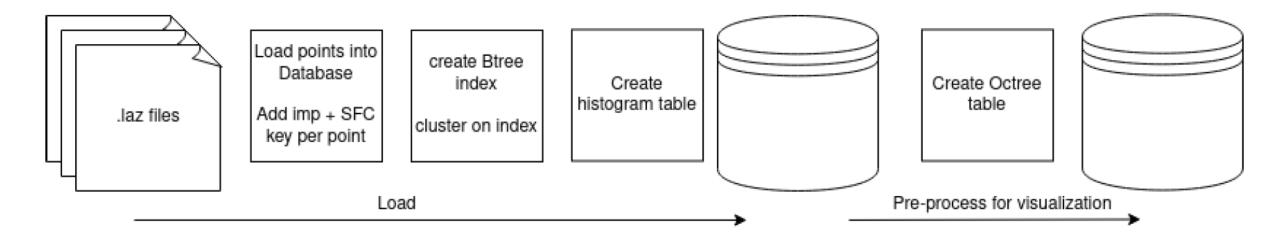
We investigate how PCServe, a web service for disseminating massive point clouds, performs for read-only access (i.e. a visualization of the control of the tion application). PCServe is backed by a database model based on Space Filling Curves. By adding a virtual hierarchy of blocks to the database, we can support different visualization applications for retrieval of point cloud data over the web without having to store the data multiple times. This makes expressive access to point clouds over the web possible. We investigate the amount of processing that is needed to create the database model and how well PCServe handles requests from the visualization application. Some suggestions are provided how the current approach can be improved.

### Implementation: PCServe



#### PCServe: DB preparation

- load, add cLoD and organize by SFC
- virtual octree to support potree viewer as there is not yet a cLol viewer
- on request generate octree nodes with SFC ranges



#### nD Convex polytope query

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journal homepage: www.elsevier.com/locate/jag



Executing convex polytope queries on nD point clouds

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#### ARTICLE INFO

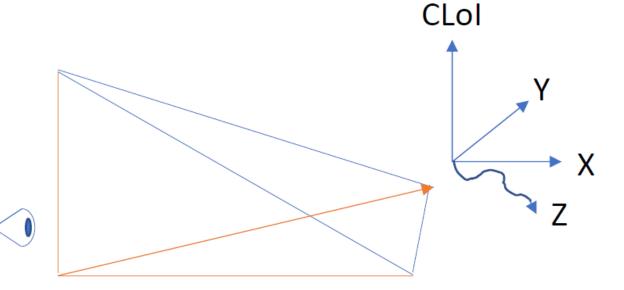
Keywords: nD point clouds Polytope query Spatial data structures CPLEX Perspective view selection

#### ABSTRACT

Efficient spatial queries are frequently needed to extract useful information from massive nD point clouds. Most previous studies focus on developing solutions for orthogonal window queries, while rarely considering the polytope query. The latter query, which includes the widely adopted polygonal query in 2D, also plays a critical role in many nD spatial applications such as the perspective view selection. Aiming for an nD solution, this paper first formulates a convex nD-polytope for querying. Then, the paper integrates three approximate geometric algorithms – SWEEP, SPHERE, VERTEX, and a linear programming method CPLEX, developing a solution based

## nD query by convex polytopes

- convex polytope is more selective than a hyper-rectangle/sphere
- well defined in nD, based on half-space intersection
- fast implementation based on SFC/virtual 2<sup>n</sup>-tree

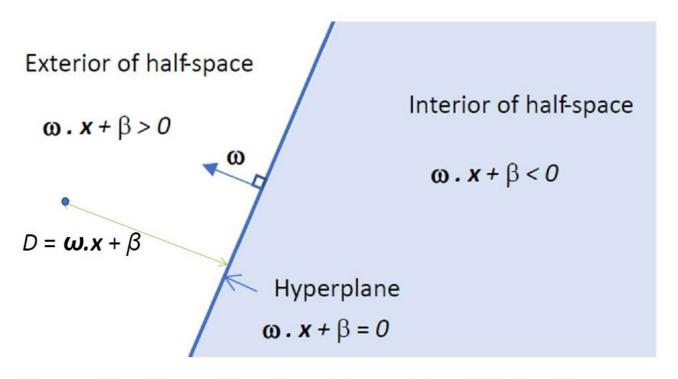


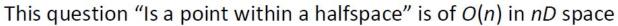
example view frustum selection

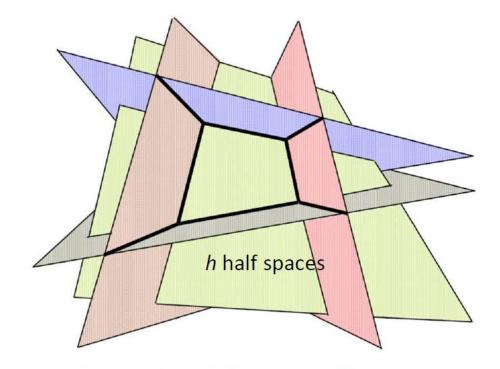
4D View Frustum with CLoI (only drawn in 3D)

#### Test overlap half-space with virtual node 2<sup>n</sup>-tree

- half-space = nD plane inequation
- virtual node 2<sup>n</sup>-tree = hyperbox (range of SFC-keys)
- plane normal vector used which 2 corners hyperbox used for test



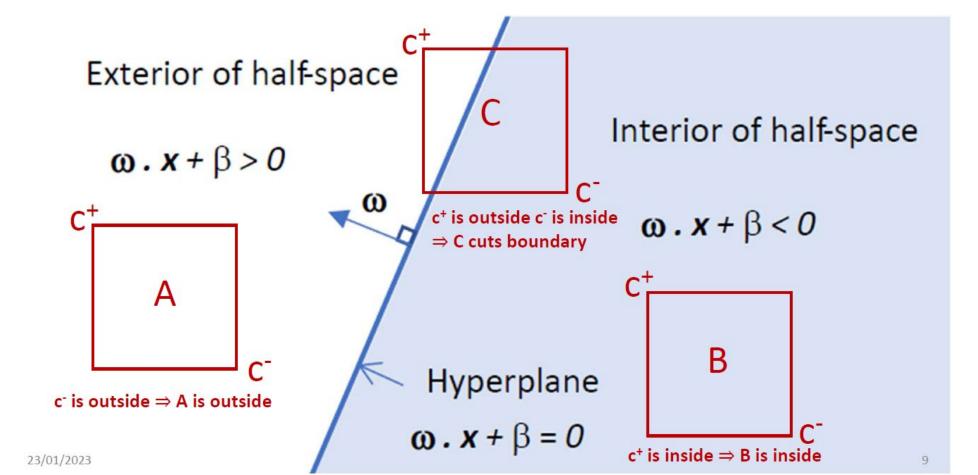




So the question of "Is a point within a region defined by h half spaces" is O(nh).

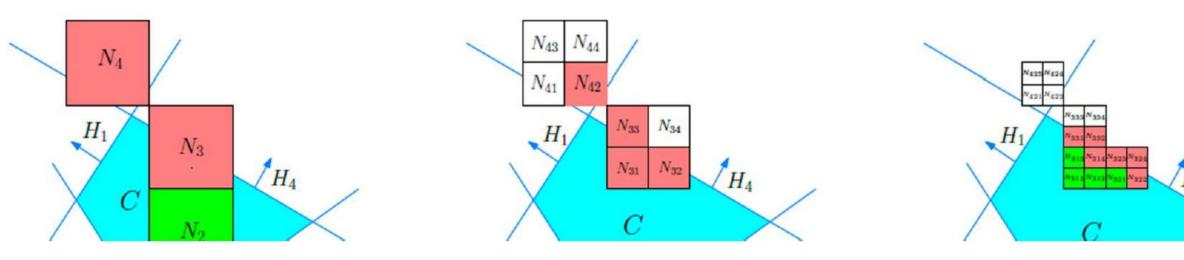
#### Just 2 corner points

- test nD box inside nD half-space: out, partial, in
- in nD box there are  $2^n$  points (e.g. 5D  $\rightarrow$  32 points), but just 2 tested



## Full query: recursive 2<sup>n</sup>-tree traversal

- in-box give SFC-range (based on level in tree this can be larger/smaller)
- partial-box is refined (unless max depth/bottom level is reached)
- resulting set of SFC ranges is joined with SFC based nD points
  - either report individual points, or point groups



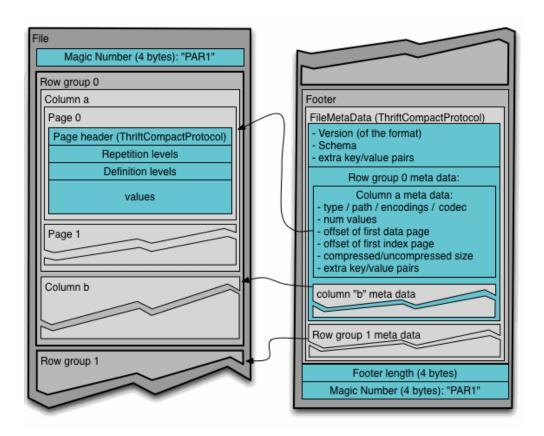
## Apache Parquet (as alternative store)

- columnar store format, for serialized data (disk storage)
- structured (+ nested) data, compressed
- has a main memory cousin (uncompressed columnar)
  - Apache Arrow → Data frame libraries can read Parquet from disk into Arrow(2D, heterogeneous tabular data)
- you can interact via SQL with the Parquet files using query engine (DuckDb, Apache Drill)



#### Apache Parquet - structure

- table split over:
  - Row groups → Columns
  - Column Chunks → Pages
- encoding data + Compression at Page level
- metadata on Min/Max of Row groups (+ recent addition: Pages), stored at end of file (block ranges index, if data sorted!)
- can be fetched in parts with http head followed by http range requests (if server support)
- C++, Java, Rust libraries read / write
- https://parquet.apache.org/



## Some nD PC data (load similar to PostgreSQL)

#### SFC key

#### **Attributes**

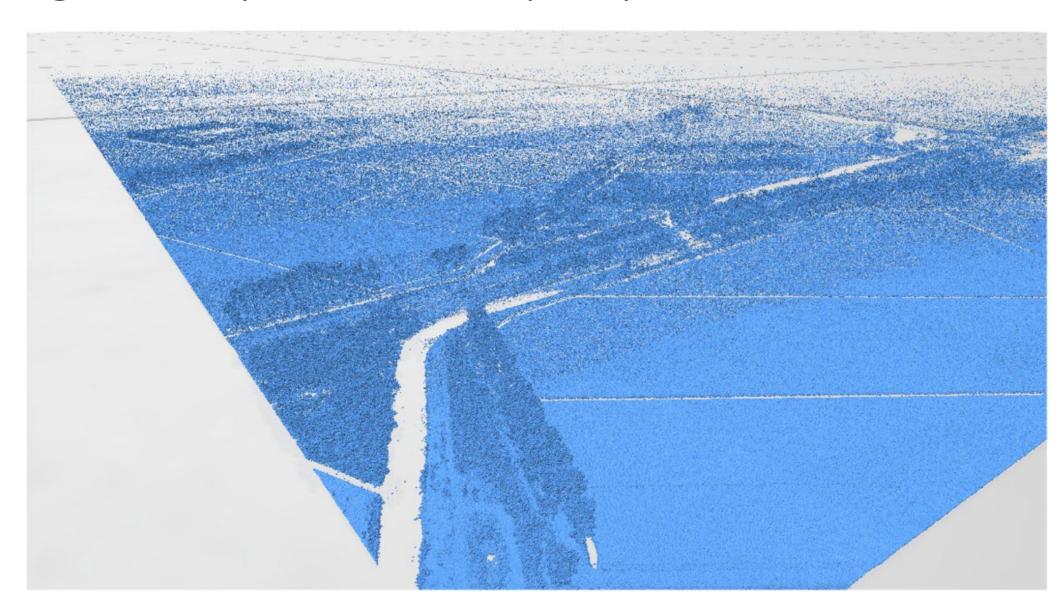
sfc_key blob	gps_time double	scan_angle float	intensity uint16	return_number uint8	number_of_returns uint8	classification uint8	scan_direction uint8	is_edge_of_flight_line uint8	
\x00\x00\x00\x00\x10\x02~\x8C\xC3,N\x8ES\xC \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFC\x04 \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFC\x04 \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFC\xC9 \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFF\xF9 \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFF\xCB \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFF\xCB \x00\x00\x00\x00\x10\x02~\x8C\xC3,N\xFF\xD7 \x00\x00\x00\x00\x10\x02~\x8C\xC3,O\xF7X\x0 \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x0E\x8C \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x0E\x8C \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x0E\x8B \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x0E\x8B \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x10\x8A} \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x10\x8A} \x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x10\x8A} \x00\x00\x00\x00\x00\x10\x02~\x8C\xC3,P\x10\x8A}	494147.3131107338 493868.6779278455 494146.86342112836 494146.83669732255 493869.0259201808 494146.41376256285 493868.97246252943 493870.1446405349 493870.0054991533 493870.1339485246 493870.0858279384 493870.03758022806 493869.855603322 493869.74327855057	27.0 16.002 27.0 27.0 16.002 28.002 15.0 16.002 15.0 16.002 15.0 16.002 15.0 16.002	170 118 22 37 64 30 184 230 66 195 48 200 47 200	1 1 1 1 1 1 1 1 1 1 1 1 1 1	1 1 1 1 1 1 1 1 1 1 1 1 1	6 2 2 2 6 6 6 6 6 6 6	1 1 1 1 1 1 1 1 1 1 1 1 1	0 0 0 0 0 0 0 0 0	
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6083217 rows (40 shown) 9 columns									(

Row groups + Columns = Column Chunks

## Querying Parquet

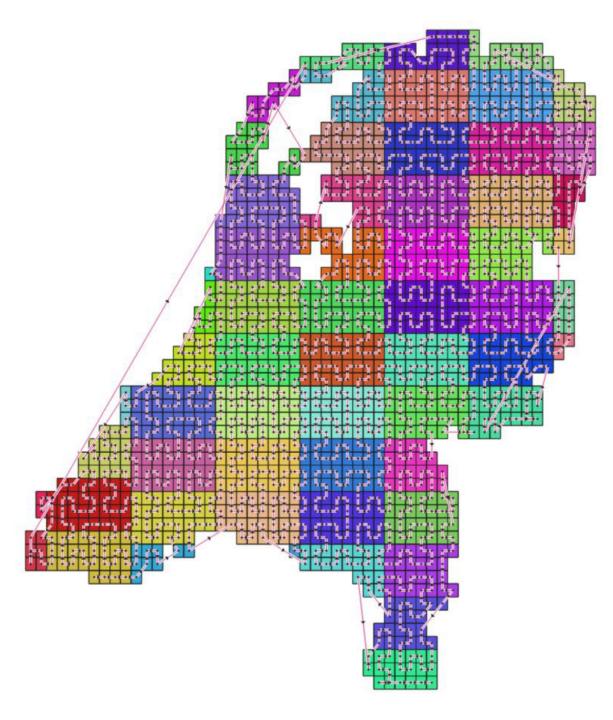
- Column Chunks divided into Pages
- encoding/compression can differ per column data type
- custom merge join with half spaces → query ranges and data (SFC keys)Read relevant parts from Parquet file, that contains sorted data
- heavy use of Min/Max index data on both Row groups as well as Column pages in footer of Parquet file
- fetch in phases:
  - 1. matching records based on SFC key column (a. Row Group  $\rightarrow$  b. Column page)
  - 2. retrieve additional attributes for query result points from disk
- decode / de-quantize keys

## Integrated Space / cLol query



## Agenda

- nD-PointCloud
  - foundations
  - PostgreSQL implementation
  - nD Convex polytope query
  - Apache Parquet
- practical results
  - Potree conversions AHN 1 to 5
  - COPC/VPC conversions
  - fast direct PC-based change detection
- future work
  - main publications



## Dell PowerEdge R7525 > ronna physical server for benchmarks, parallel computing

#### Dell PowerEdge R7525 server

- 1. 2 x 24-core AMD EPYC 7443 processors (96 threads), 2.85GHz
- 2. 512 GB buffered memory (RDIMM)
- 3. Ubuntu operating system

- Disk storage direct attached
  - 1. 15,4 Tb SSD SAS
  - 2. 19,2 Tb SSD SATA
  - 3. 180,0 Tb Hard Drive SAS, 7.2K rpm in RAID-5

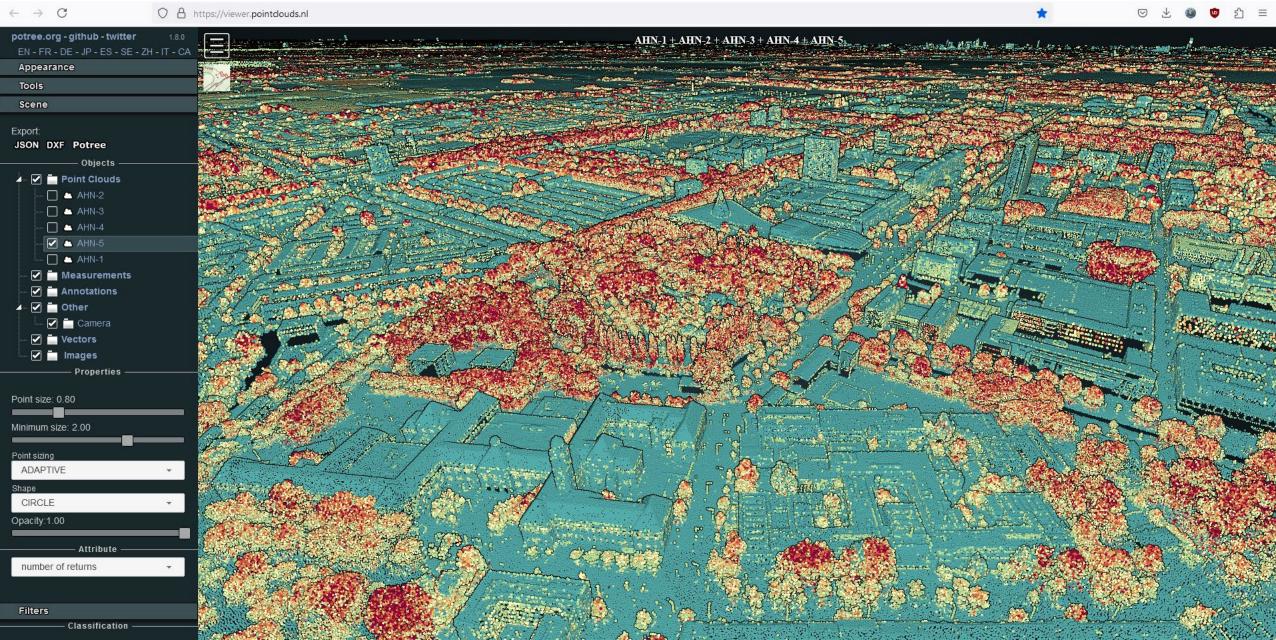


Network NAS 26,0 TB on svm104.storage.tudelft.net

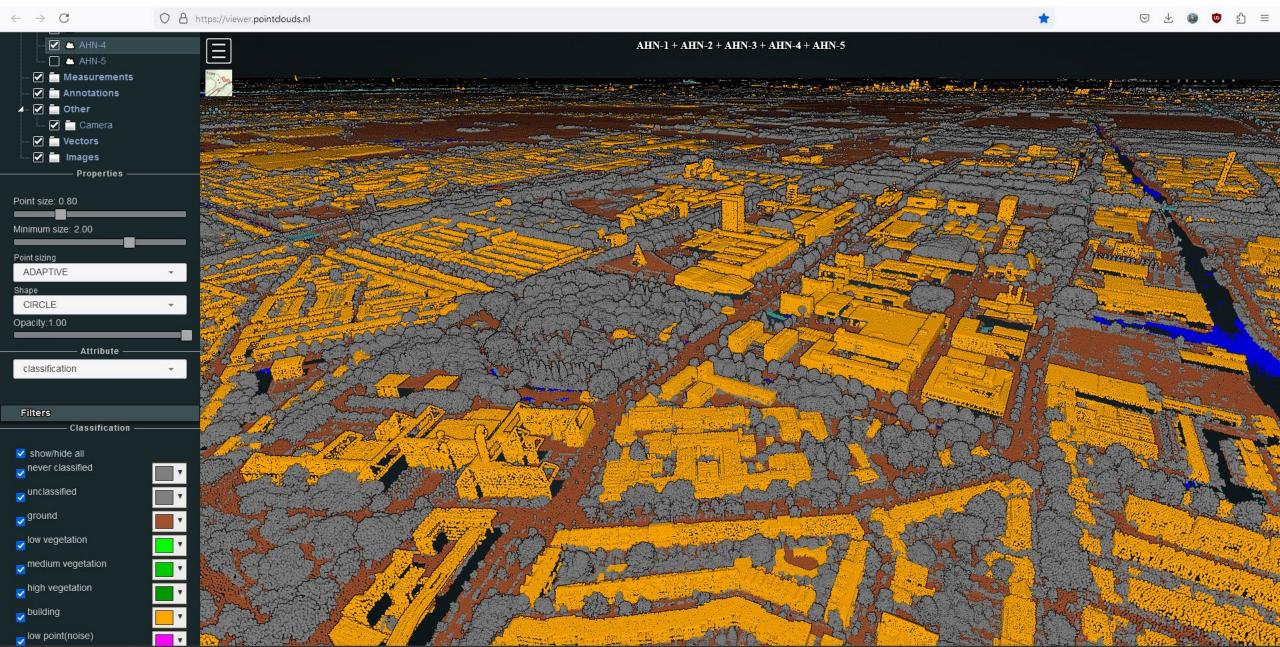
### http://viewer.pointclouds.nl all AHN versions in potree



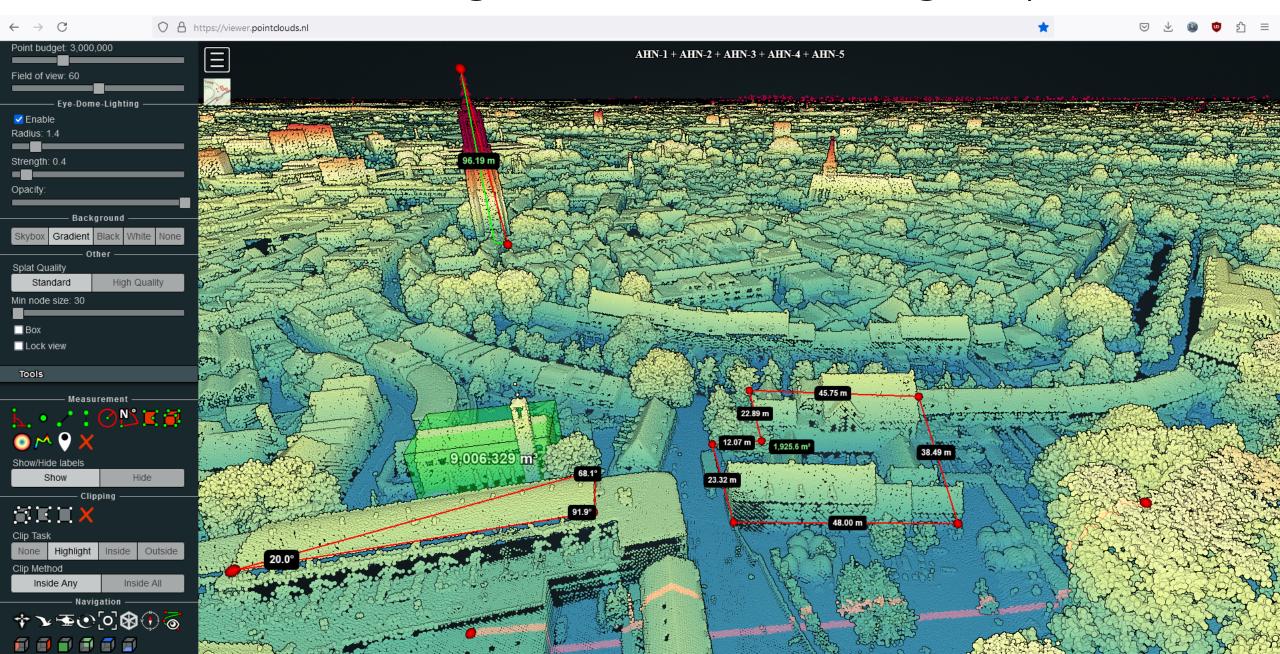
All potree tools, e.g. show number of returns: 1-7



### .. or classification



#### Various tools: height, area, volume, angles, profile,...

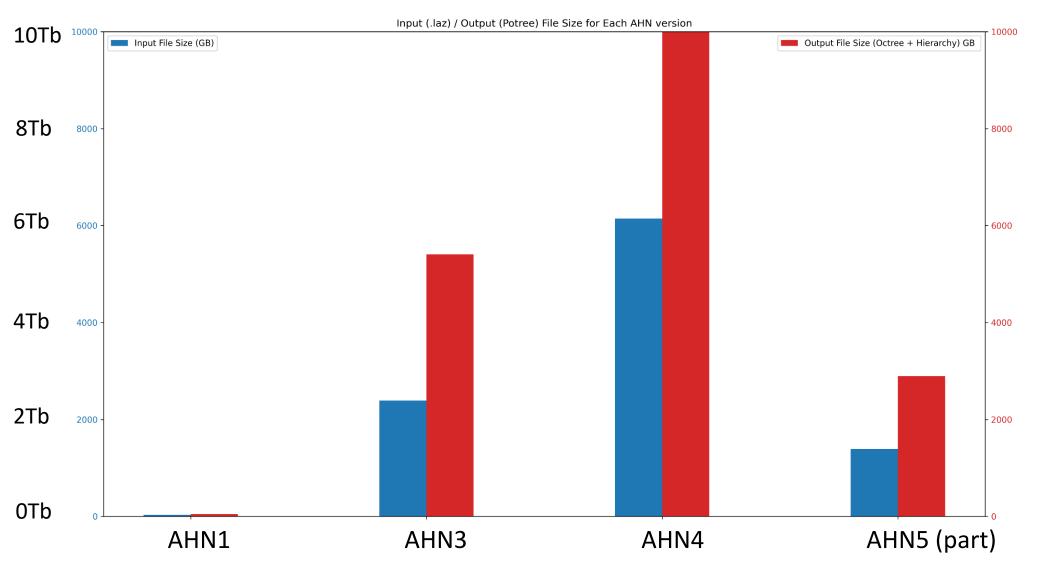


#### potree conversion on ronna

name	#points	#files in	Input size (Gb)	Output size (Gb)	Duration (hours)
AHN1	11.984.853.767	1358	33,1	44,2	0,17
AHN3	557.925.797.136	1374	2390,7	5035,8	12,29
AHN4	947.364.043.509	1381	6145,4	9931,8	33,38
AHN5 part	289.944.615.278	499	1394,3	2695,9	8,59

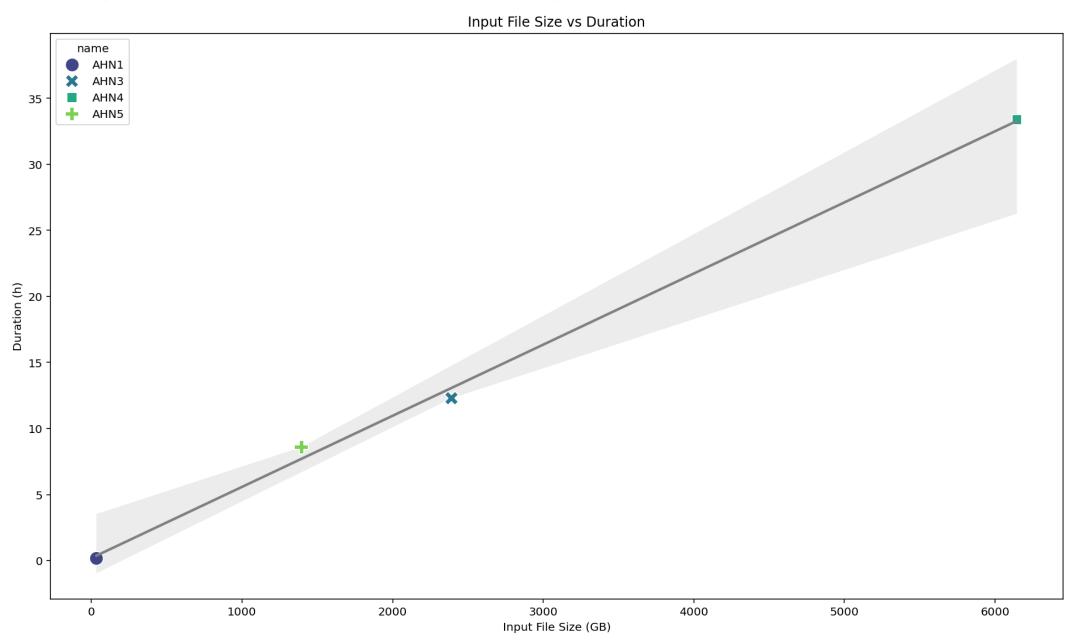
AHN2 converted before (previous project) Duration conversion on ronna, multi-user

#### AHN-potree conversion statistics: in-/output size



Note about issue memory leak potree, due to parallel sort in library (in some systems/versions)

#### AHN-potree conversion: input size vs duration



## Virtual Point Clouds and Cloud Optimized Point Cloud

- Virtual Point Cloud (VPC) file = SpatioTemporal Asset Catalogs (STAC)
- extension on top of GeoJSON with PC metadat
- contains links to COPC .laz files (or EPT, Entwine Point Tile)

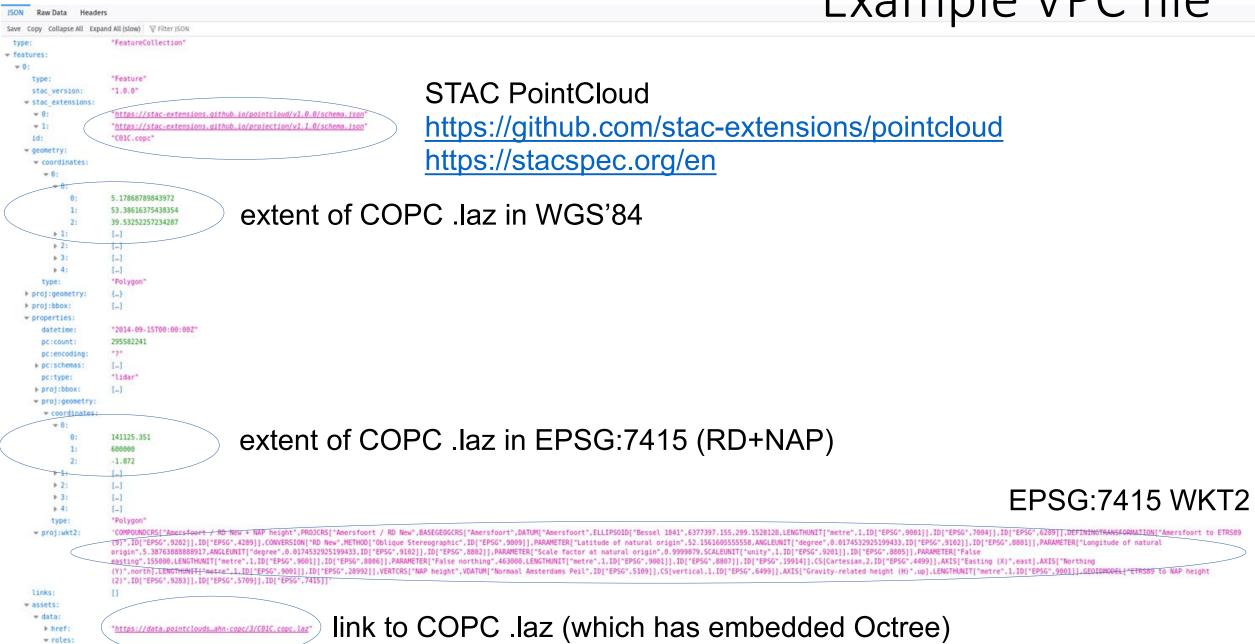


- Cloud Optimized Point Cloud (COPC) file = data
- COPC is .laz which has embedded Octree
- specs on <a href="https://copc.io/">https://copc.io/</a> (by Hobu, Inc.)



 load VPC in QGIS as a single point cloud layer (rather than each file as a separate map layer)

#### Example VPC file

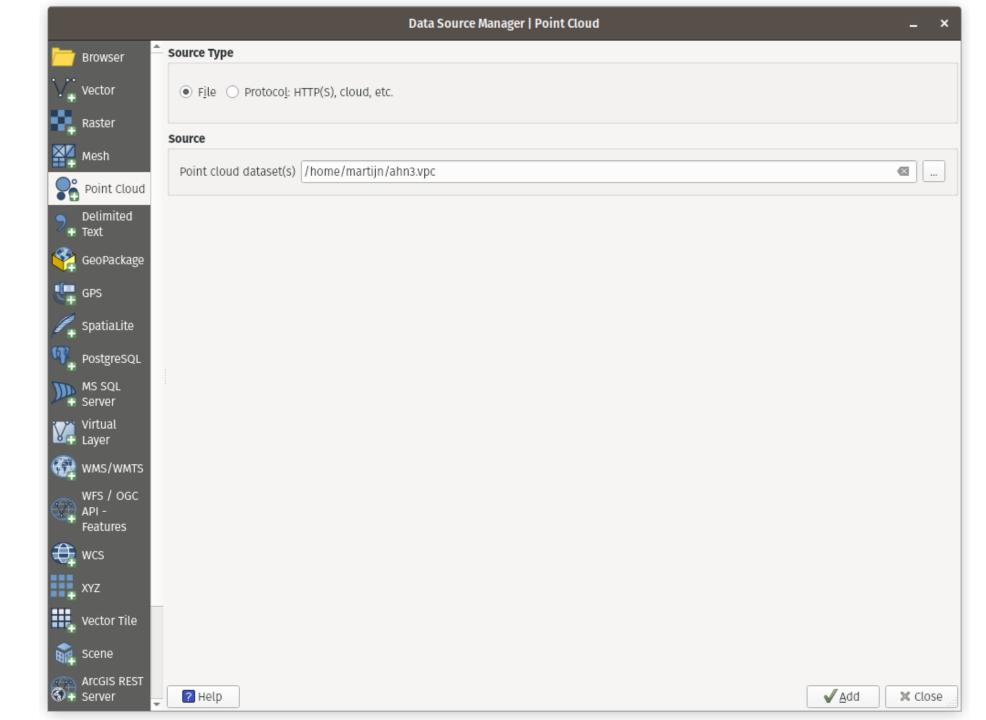


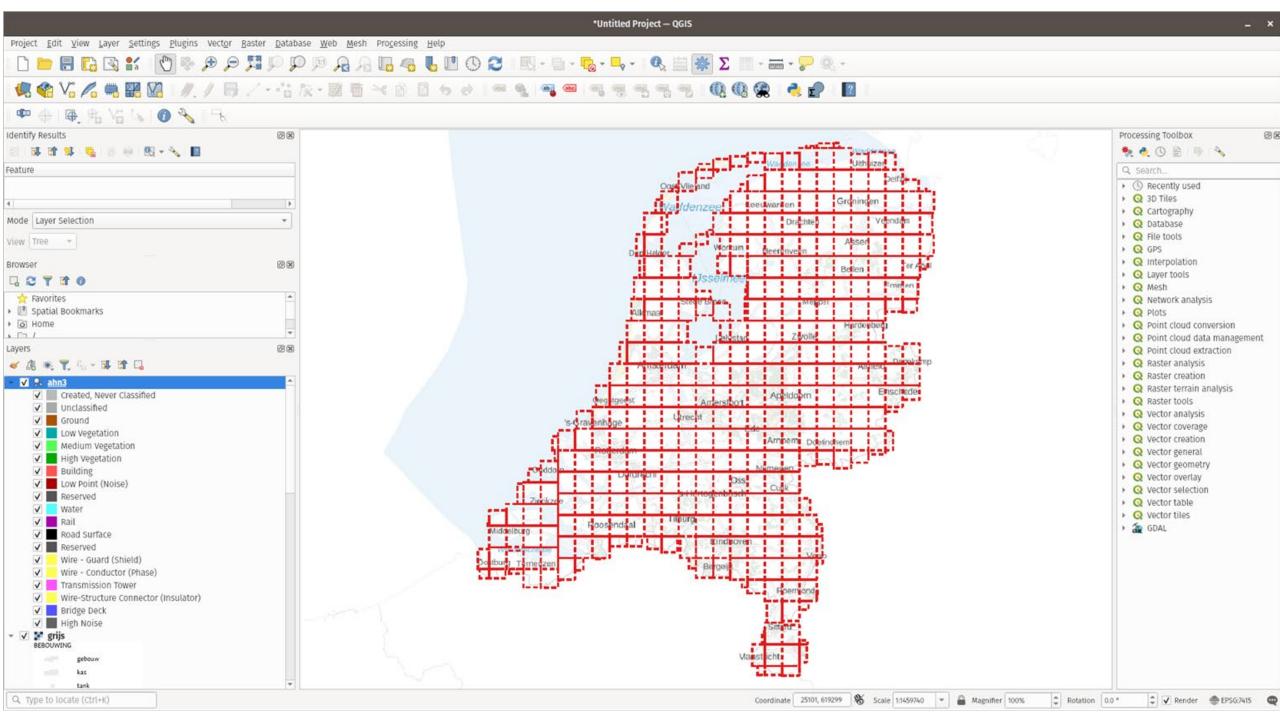
## To use VPC/COPC in QGIS start with download .vpc file

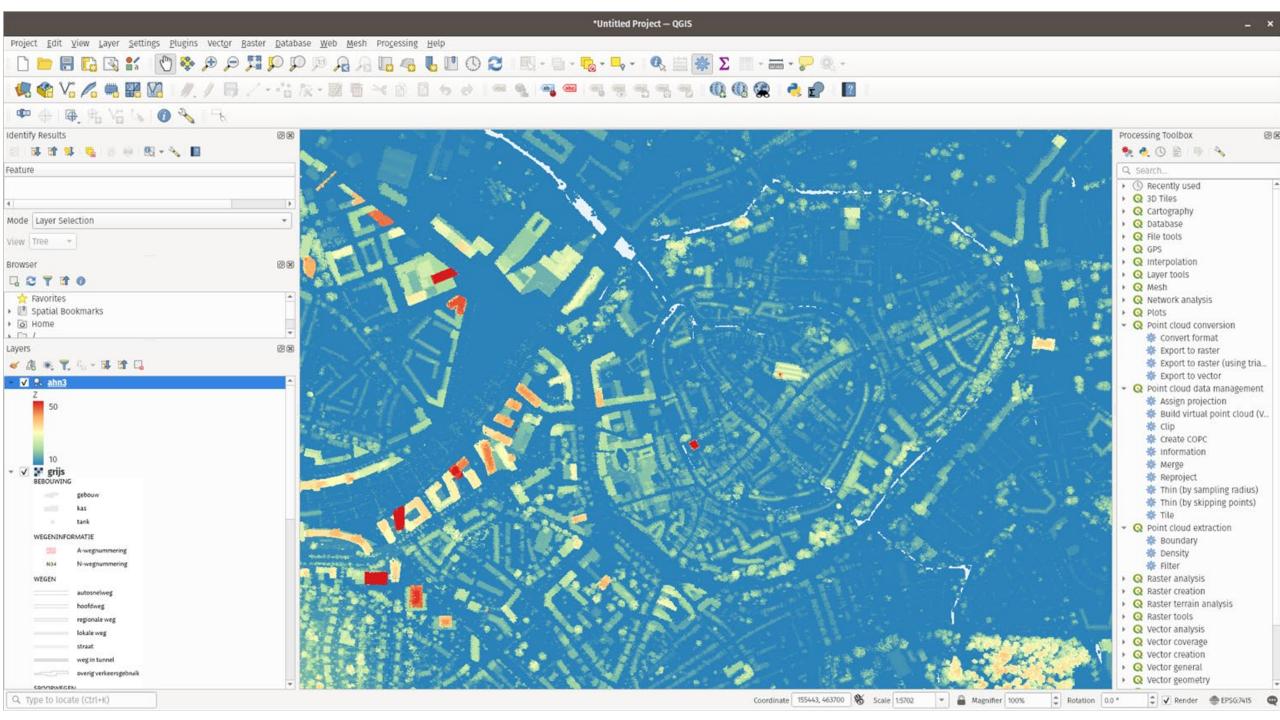
2024-10-10 14:03:00 (109 MB/s) - 'ahn3.vpc' saved [2428491/2428491]

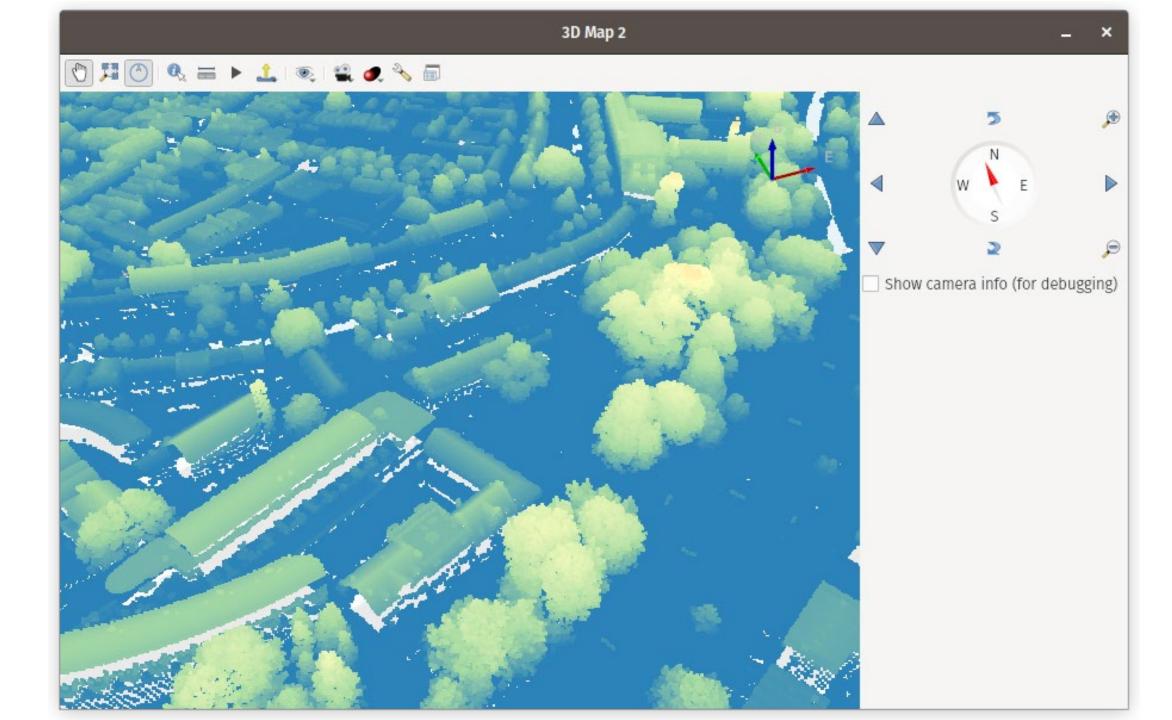
```
$ wget https://data.pointclouds.nl/data/nl/ahn-copc/3/ahn3.copc.vpc.json -O ahn3.vpc
--2024-10-10 14:03:00-- https://data.pointclouds.nl/data/nl/ahn-copc/3/ahn3.copc.vpc.json
Resolving data.pointclouds.nl (data.pointclouds.nl)... 131.180.126.49
Connecting to data.pointclouds.nl (data.pointclouds.nl) | 131.180.126.49 | : 443... connected.
HTTP request sent, awaiting response... 200 OK
Length: 2428491 (2.3M) [application/json]
Saving to: 'ahn3.vpc'
ahn3.vpc
               2.32M --.-KB/s in 0.02s
```

## Open VPC in QGIS









# Statistics for AHN to COPC and VPC (on ronna)

	#files in	input	processing	#files out	output
AHN3	1374	2.4Tb	14 days 5 hrs	382	3.1Tb
AHN4	1381	6.1TB	26 days *)	381	5.1Tb

<sup>\*)</sup> estimation (as process was interrupted, total 54 days)

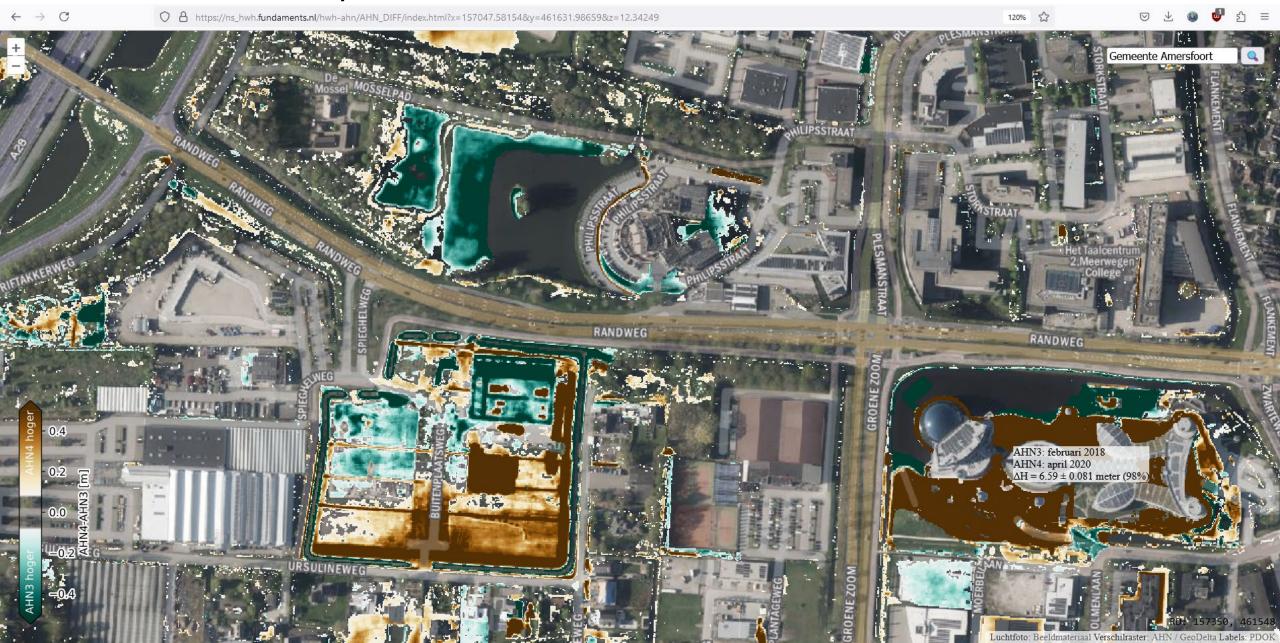
processing time linear with input size ronna is multi-user system, so timings not exact for AHN3: COPC > LAZ input, for AHN4: COPC < LAZ input

GitHub to source <a href="https://github.com/nd-pc/process">https://github.com/nd-pc/process</a> ahn

# Raster based change detection



# Nicer example AFAS HQ, Leusden



# Fine, but...

- grid is less detailed than PC
- grid can only be displayed in 2D
- grid needs to be computed first

 much better to have direct point cloud change detection

#### Comparison of Cloud-to-Cloud Distance Calculation Methods - Is the Most Complex Always the Most Suitable?



Vitali Diaz, Peter van Oosterom, Martijn Meijers, Edward Verbree, Nauman Ahmed, and Thijs van Lankveld

Abstract Cloud-to-cloud (C2C) distance calculations are frequently performed as an initial stage in change detection and spatiotemporal analysis with point clouds. There are various methods for calculating C2C distance, also called interpoint distance, which refers to the distance between two corresponding point clouds captured at different epochs. These methods can be classified from simple to complex, with more steps and calculations required for the latter. Generally, it is assumed that a more complex method will result in a more precise calculation of inter-point distance, but this assumption is rarely evaluated. This paper compares eight commonly used methods for calculating the inter-point distance. The results indicate that the accuracy of distance calculations depends on the chosen method and a characteristic related to the point density, the intra-point distance, which refers to the distance between points within the same point cloud. The results are helpful for applications that analyze spatiotemporal point clouds for change detection. The findings will be helpful in future applications, including analyzing spatiotemporal point clouds for change detection.

Keywords Cloud-to-cloud distance calculation · Change detection · Spatiotemporal analysis

This article was selected based on the results of a double-blind review of an extended abstract

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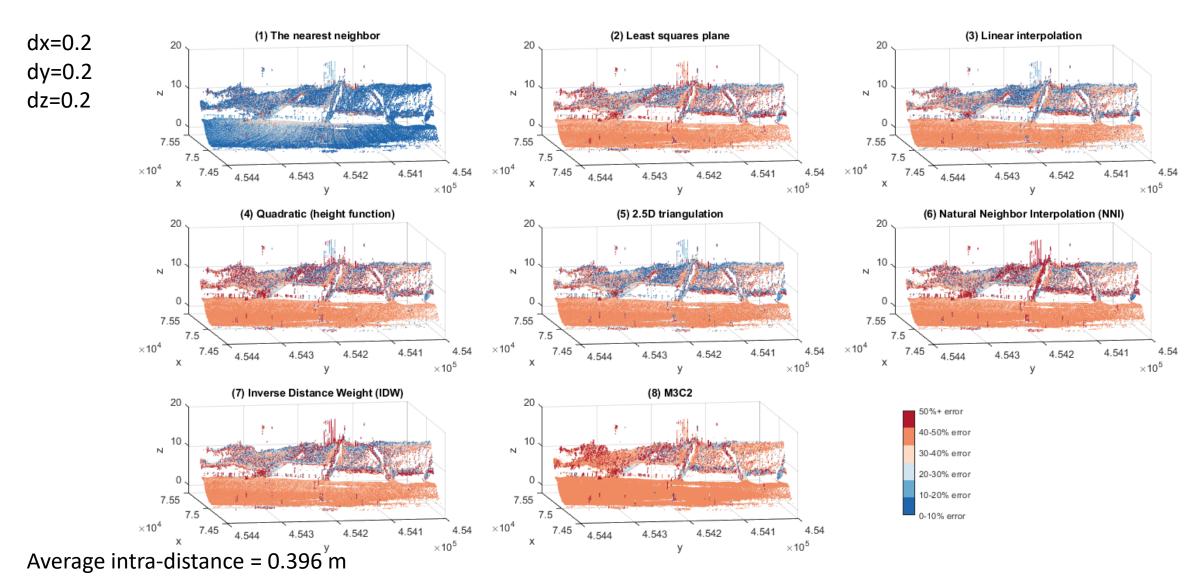
© The Author(s), under exclusive license to Springer Nature Switzerland AG 2024 T. H. Kolbe et al. (eds.), Recent Advances in 3D Geoinformation Science, Lecture Notes in Geoinformation and Cartography, https://doi.org/10.1007/978-3-031-43699-4\_20

# Direct PC-based change detection (not convert to grids and show changes in 3D)

- tests to find the suitable PC-based method: controlled offsets and real epochs
   comparison based on error and time cost
- range of algorithms (as available in CloudCompared)
- reimplemented in Matlab for full control
- testing with various data sets (AHN, CostScan, lake, bunny)
- conclusion from results: 3D Nearest Neighbour(NN) outperforms/equal quality to more complex algorithms regarding similar results and less time cost. Relatively efficient for implementation, 'real' NN kd-tree based.
- Then, further speed-up will be challenged for massive implementation for further change detection at nationwide applications over the AHN series.

## Experiment 1: CoastScan data with controlled offset

comparing the 8 different algorithm



## Experiment 2: Preliminary Results (2D NN)

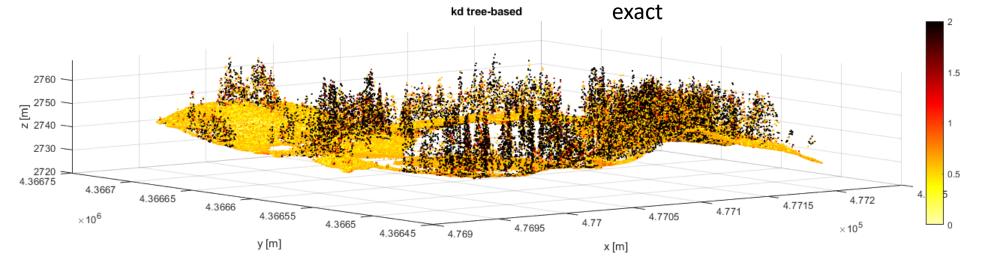
#### Artificial displacement

dx = -0.25

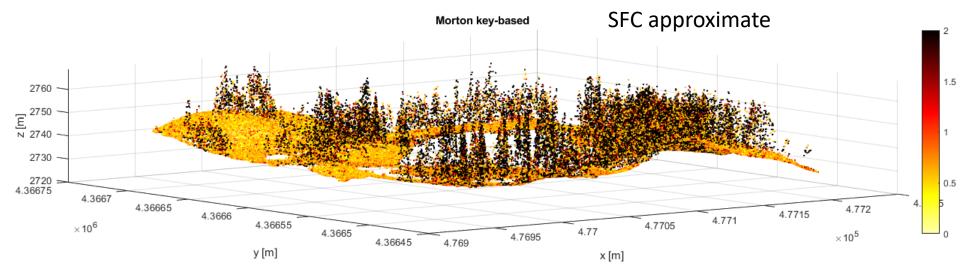
dy = -0.25

dz = -0.5

d=0.6124



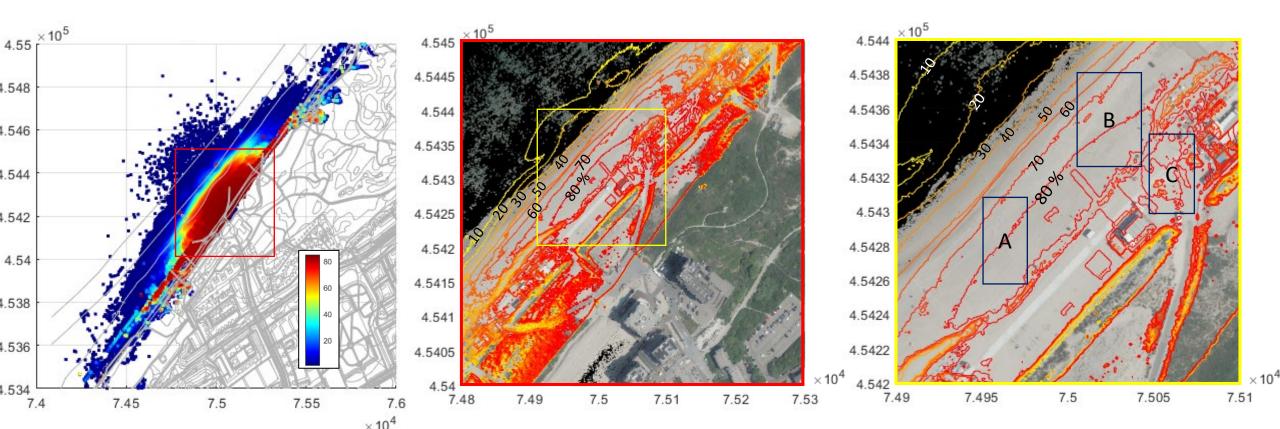
(lesson learnt: do not do 2D, but 3D NN)



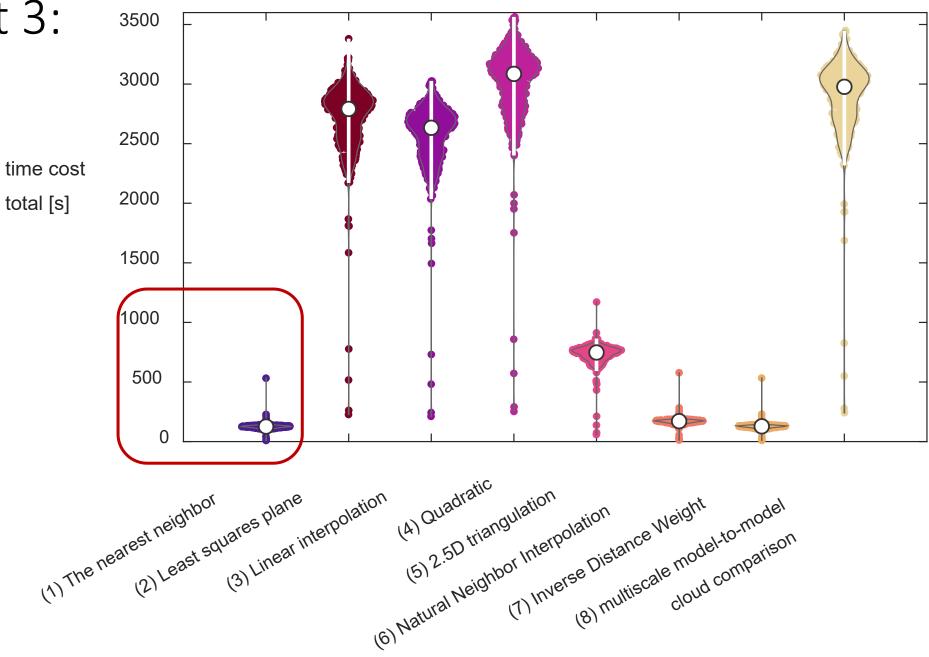
### Experiment 3: spatio-temporal CoastScan data



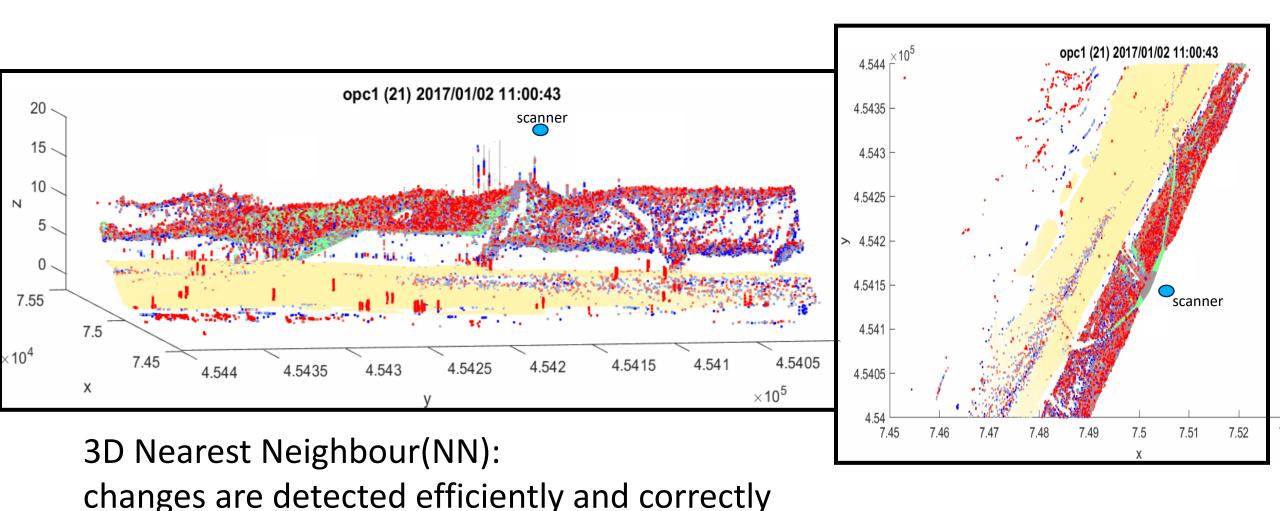
period 2017/01/01 to 2017/04/18
2,187 epochs



Experiment 3: Time cost



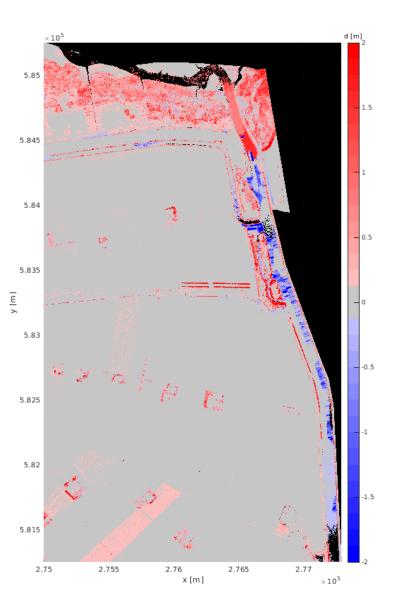
# Experiment 3: Spatio-temporal CoastScan data



# SFC-based ultra-fast change detection

- per epoch, preparation of PC data set
  - 1. scaling and offsetting (SFC always in a cube)
  - 2. SFC key calculation
  - 3. sorting SFC key column
- batch process NN distance calculation (advanced merge join 2 epochs): for every point P in epoch 2:
  - 1. search the previous and next SFC key of the calculated SFC key in epoch 1
  - 2. decode and unscale these two SFC keys (previous and next)
  - 3. find approximate nearest neighbor (NN) of the 2 candidates (option: box query in epoch 1 with P as center and distance P-NN as radius)
  - 4. output d and dx,dy,dz

# Preliminary results: AHN4-3





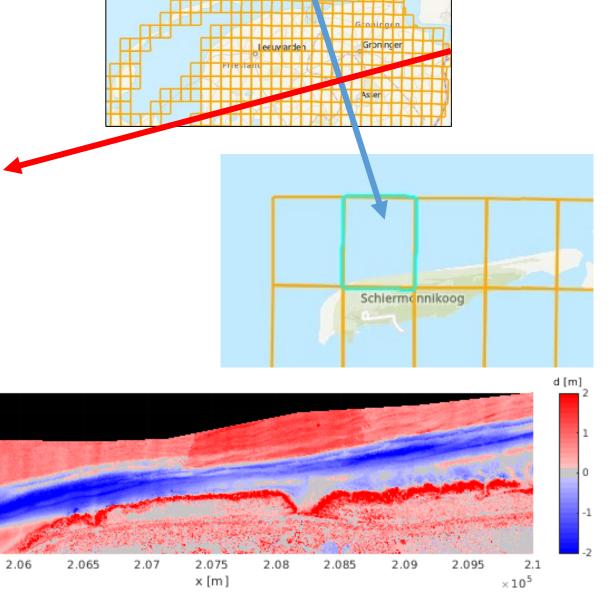
6.136 6.134

E 6.132 > 6.13

6.128

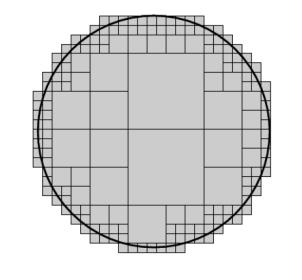
2.05

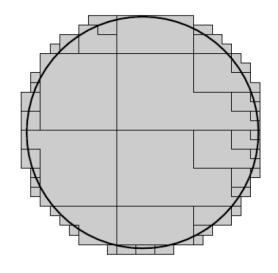
2.055

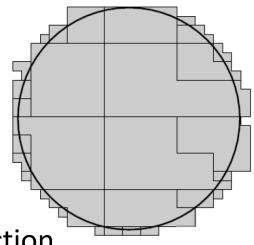


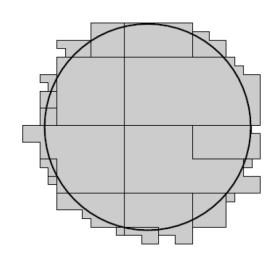
# Agenda

- nD-PointCloud
  - foundations
  - PostgreSQL implementation
  - nD Convex polytope query
  - Apache Parquet
- practical results
  - Potree conversions AHN 1 to 5
  - COPC/VPC conversions
  - fast direct point cloud-based change detection









- future work
  - main publications

#### Conclusion

- nD-PointCloud
- designed and tested nD-PC organization (with cLoI) in resp. DBMS and files, and good results obtained
- explored and cleaned spatio-temporal point cloud data, and analyzed in detail the possible change detection options → NN preferred
- current practice tools not yet cLol-aware and space first (not time)
- active at various meetings (OGC, 3D GeoInfo, FOSS4G, EGU) and publications
- Geomatics/GIMA students doing (thesis) projects with point clouds
- open invitation to the nD-PointCloud project concluding symposium:
   12 November 2024, 13-16 hours
   room B, Faculty of Architecture and the Built Environment, TU Delft

#### Future work

- completion of publication in pipeline
- standardization of format and protocols (binary Parquet files)
- cLoI aware viewers and other clients (computations), using selection with flat hyperplanes (convex polytope)
- on-the-fly CRS transformations of selections, or pre-computing of whole data sets (on HPC/HTC)
- explore cLoI to integrate datasets from different scales (after georeferencing/CRS transformation)
- nationwide AHN change detection (buildings, vegetation, ..) by
  - adding 1 just or all 3 directions
  - store result integrated (dimension of attribute) or separate from points
  - interactive of massive preprocessing (on HPC/HTC), but fast use
  - backward of forward changes
  - for CoastScan, option to skip epochs (hours, days, weeks, months)

# Publications (1/2)

- Vitali Diaz, Peter van Oosterom, Martijn Meijers, Edward Verbree, Nauman Ahmed, Thijs van Lankveld, Comparison of Cloud-to-Cloud Distance Calculation Methods - Is the Most Complex Always the Most Suitable?, Chapter in: Recent Advances in 3D Geoinformation Science, Lecture Notes in Geoinformation and Cartography, Springer Nature Switzerland, pp. 229-334, 2024.
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- Zhenyu Liu, Peter van Oosterom, Jesús Balado, Arjen Swart, Bart Beers, Data frame aware optimized Octomap-based dynamic object detection and removal in Mobile Laser Scanning data, In: Alexandria Engineering Journal, 74, pp. 327-344, 2023.
- Vitali Diaz, Peter van Oosterom, Martijn Meijers, Edward Verbree, Nauman Ahmed, Thijs van Lankveld, Comparison of point distance calculation methods in point clouds - Is the most complex always the most suitable?, In: Proceedings of the 18th International 3DGeoInfo Conference 2023, Munich, Germany, pp. 329-334, 2023.
- Ioannis Dardavesis, Edward Verbree, Azarakhsh Rafiee, Indoor localisation and location tracking in indoor facilities based on LiDAR point clouds and images of the ceilings, In: Proceedings of the 26th AGILE Conference on Geographic Information Science, 2023, GIScience Series, 4(4), Delft, The Netherlands, pp. 1-15, 2023
- Vidushi Bhatt, Sharath Chandra Madanu, Shen Qiwei, Susanne Epema, Gees Brouwer, Pointcloud based anatomy, MSc Geomatics synthesis project, Technical report, Delft University of Technology, pp. 55, 2023.
- Haicheng Liu, nD-PointCloud Data Management continuous levels, adaptive histograms, and diverse
  query geometries, PhD thesis, Delft University of Technology, pp. 207, 2022.

# Publications (2/2)

- Peter van Oosterom, Simon van Oosterom, Haicheng Liu, Rod Thompson, Martijn Meijers, Edward Verbree, Organizing and visualizing point clouds with continuous levels of detail, In: ISPRS Journal of Photogrammetry and Remote Sensing, Elsevier BV, 194, pp. 119—131, 2022.
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- Vitali Diaz, Haicheng Liu, Peter van Oosterom, Martijn Meijers, Edward Verbree, Fedor Baart, Maarten Pronk, Thijs van Lankveld, Point clouds and Hydroinformatics, 2022 (Abstract from EGU General Assembly 2022, Vienna, Austria, 23–27 May 2022).
- Zhenyu Liu, Peter van Oosterom, Jesús Balado, Arjen Swart, Bart Beers, Detection and reconstruction of static vehicle-related ground occlusions in point clouds from mobile laser scanning, In: Automation in Construction, Elsevier BV, 141, pp. 104461, 2022.
- Haicheng Liu, Rodney Thompson, Peter van Oosterom, Martijn Meijers, Executing convex polytope queries on nD point clouds, In: International Journal of Applied Earth Observations and Geoinformation, Elsevier, 105(102625), pp. 1-11, 2021.
- Guan-Ting Zhang, Edward Verbree, Xiao-Jun Wang, An Approach to Map Visibility in the Built Environment From Airborne LiDAR Point Clouds, In: IEEE Access, Institute of Electrical and Electronics Engineers (IEEE), 9, pp. 44150-44161, 2021.
- H. Liu, P. Van Oosterom, B. Mao, M. Meijers, R. Thompson, An efficient nD-Point Data Structure for Querying Flood Risks, Chapter in: The International Archives of the Photogrammetry, Remote Sensing and Spatial Information Sciences, Copernicus GmbH, XLIII-B4-2021, pp. 367-374, 2021.

Thanks for your attention!



